

# Math 571 Chapter 2 Multiplicative Structures

Robert C. Vaughan

January 6, 2023

- In elementary number theory courses it is usual taught that the reduced residue classes modulo  $q$  form a cyclic group under multiplication if and only if  $q = p^k$  with  $p = 2$  and  $k = 1$  or  $2$ , or with  $p > 2$  and all  $k \geq 1$ . A generator  $g$  is called a primitive root. It is often also shown that if  $p = 2$  and  $k \geq 3$ , then every reduced residue modulo  $2^k$  is generated by

$$(-1)^u 5^v$$

where  $u = 0$  or  $1$  and  $0 \leq v < 2^{k-2}$ .

- In elementary number theory courses it is usual taught that the reduced residue classes modulo  $q$  form a cyclic group under multiplication if and only if  $q = p^k$  with  $p = 2$  and  $k = 1$  or  $2$ , or with  $p > 2$  and all  $k \geq 1$ . A generator  $g$  is called a primitive root. It is often also shown that if  $p = 2$  and  $k \geq 3$ , then every reduced residue modulo  $2^k$  is generated by

$$(-1)^u 5^v$$

where  $u = 0$  or  $1$  and  $0 \leq v < 2^{k-2}$ .

- One can then use the Chinese Remainder Theorem to express each residue modulo  $q$  in a suitable form. This was all first proved by Gauss.

- In elementary number theory courses it is usual taught that the reduced residue classes modulo  $q$  form a cyclic group under multiplication if and only if  $q = p^k$  with  $p = 2$  and  $k = 1$  or  $2$ , or with  $p > 2$  and all  $k \geq 1$ . A generator  $g$  is called a primitive root. It is often also shown that if  $p = 2$  and  $k \geq 3$ , then every reduced residue modulo  $2^k$  is generated by

$$(-1)^u 5^v$$

where  $u = 0$  or  $1$  and  $0 \leq v < 2^{k-2}$ .

- One can then use the Chinese Remainder Theorem to express each residue modulo  $q$  in a suitable form. This was all first proved by Gauss.
- It is also an example of the theorem, usually proved in abstract algebra courses, that each abelian group is a direct product of cyclic groups. The methods of abstract algebra do not necessarily give explicit representations, which are sometimes the easiest way of seeing things.

- There is a more abstract and general treat of characters which I have put in the files section if you are interested.

- There is a more abstract and general treat of characters which I have put in the files section if you are interested.
- A Dirichlet character is an arithmetical function  $\chi : \mathbb{N} \rightarrow \mathbb{C}$  with the following properties.

- There is a more abstract and general treat of characters which I have put in the files section if you are interested.
- A Dirichlet character is an arithmetical function  $\chi : \mathbb{N} \rightarrow \mathbb{C}$  with the following properties.
- 1.  $\chi$  is totally multiplicative.

- There is a more abstract and general treat of characters which I have put in the files section if you are interested.
- A Dirichlet character is an arithmetical function  $\chi : \mathbb{N} \rightarrow \mathbb{C}$  with the following properties.
  1.  $\chi$  is totally multiplicative.
  2.  $\chi$  has period  $q$  for some  $q \in \mathbb{N}$ .



- There is a more abstract and general treat of characters which I have put in the files section if you are interested.
- A Dirichlet character is an arithmetical function  $\chi : \mathbb{N} \rightarrow \mathbb{C}$  with the following properties.
  - 1.  $\chi$  is totally multiplicative.
  - 2.  $\chi$  has period  $q$  for some  $q \in \mathbb{N}$ .
  - 3. If  $(x, q) > 1$ , then  $\chi(x) = 0$ .

- There is a more abstract and general treat of characters which I have put in the files section if you are interested.
- A Dirichlet character is an arithmetical function  $\chi : \mathbb{N} \rightarrow \mathbb{C}$  with the following properties.
  - 1.  $\chi$  is totally multiplicative.
  - 2.  $\chi$  has period  $q$  for some  $q \in \mathbb{N}$ .
  - 3. If  $(x, q) > 1$ , then  $\chi(x) = 0$ .
- In view of the periodicity we can immediately extend the definition to  $\mathbb{Z}$ .

- There is a more abstract and general treat of characters which I have put in the files section if you are interested.
- A Dirichlet character is an arithmetical function  $\chi : \mathbb{N} \rightarrow \mathbb{C}$  with the following properties.
  - 1.  $\chi$  is totally multiplicative.
  - 2.  $\chi$  has period  $q$  for some  $q \in \mathbb{N}$ .
  - 3. If  $(x, q) > 1$ , then  $\chi(x) = 0$ .
- In view of the periodicity we can immediately extend the definition to  $\mathbb{Z}$ .
- From the theory of multiplicative functions we have  $\chi(1) = 1$ .

- There is a more abstract and general treat of characters which I have put in the files section if you are interested.
- A Dirichlet character is an arithmetical function  $\chi : \mathbb{N} \rightarrow \mathbb{C}$  with the following properties.
  1.  $\chi$  is totally multiplicative.
  2.  $\chi$  has period  $q$  for some  $q \in \mathbb{N}$ .
  3. If  $(x, q) > 1$ , then  $\chi(x) = 0$ .
- In view of the periodicity we can immediately extend the definition to  $\mathbb{Z}$ .
- From the theory of multiplicative functions we have  $\chi(1) = 1$ .
- The special character which is 1 whenever  $(x, q) = 1$  is called the principal character and is often denoted by  $\chi_0$ .

- By Fermat-Euler, when  $(x, q) = 1$  we have

$$1 = \chi(1) = \chi(x^{\phi(q)}) = \chi(x)^{\phi(q)},$$

- By Fermat-Euler, when  $(x, q) = 1$  we have

$$1 = \chi(1) = \chi(x^{\phi(q)}) = \chi(x)^{\phi(q)},$$

- so  $\chi(x)$  is a  $\phi(q)$ -th root of unity.

- By Fermat-Euler, when  $(x, q) = 1$  we have

$$1 = \chi(1) = \chi(x^{\phi(q)}) = \chi(x)^{\phi(q)},$$

- so  $\chi(x)$  is a  $\phi(q)$ -th root of unity.
- Also  $|\chi(x)| = 1$ .

- By Fermat-Euler, when  $(x, q) = 1$  we have

$$1 = \chi(1) = \chi(x^{\phi(q)}) = \chi(x)^{\phi(q)},$$

- so  $\chi(x)$  is a  $\phi(q)$ -th root of unity.
- Also  $|\chi(x)| = 1$ .
- Hence the number of possible characters modulo  $q$  is at most  $\phi(q)^{\phi(q)}$ , i.e. is finite.



- By Fermat-Euler, when  $(x, q) = 1$  we have

$$1 = \chi(1) = \chi(x^{\phi(q)}) = \chi(x)^{\phi(q)},$$

- so  $\chi(x)$  is a  $\phi(q)$ -th root of unity.
- Also  $|\chi(x)| = 1$ .
- Hence the number of possible characters modulo  $q$  is at most  $\phi(q)^{\phi(q)}$ , i.e. is finite.
- Let their number be  $h$ .

- If  $(a, q) = 1$ , then

$$\sum_{x=1}^q \chi(x) = \sum_{x=1}^q \chi(ax) = \chi(a) \sum_{x=1}^q \chi(x)$$

- If  $(a, q) = 1$ , then

$$\sum_{x=1}^q \chi(x) = \sum_{x=1}^q \chi(ax) = \chi(a) \sum_{x=1}^q \chi(x)$$

- Hence if there is an  $a$  with  $(a, q) = 1$  and  $\chi(a) \neq 1$ , then the sum is 0.

- If  $(a, q) = 1$ , then

$$\sum_{x=1}^q \chi(x) = \sum_{x=1}^q \chi(ax) = \chi(a) \sum_{x=1}^q \chi(x)$$

- Hence if there is an  $a$  with  $(a, q) = 1$  and  $\chi(a) \neq 1$ , then the sum is 0.
- Thus we have

### Lemma 1

*Suppose that  $\chi$  is a character modulo  $q$ . Then*

$$\frac{1}{\phi(q)} \sum_{x=1}^q \chi(x) = \begin{cases} 1 & (\chi = \chi_0) \\ 0 & (\chi \neq \chi_0). \end{cases}$$

- If  $\chi_1$  and  $\chi_2$  are characters modulo  $q_1$  and  $q_2$  respectively, then  $\chi_1\chi_2$  is one modulo  $q_1q_2$ .

- If  $\chi_1$  and  $\chi_2$  are characters modulo  $q_1$  and  $q_2$  respectively, then  $\chi_1\chi_2$  is one modulo  $q_1q_2$ .
- If  $\chi$  is a character, then so is  $\bar{\chi}$ , and  $\chi\bar{\chi} = \bar{\chi}\chi = \chi_0$ .

- If  $\chi_1$  and  $\chi_2$  are characters modulo  $q_1$  and  $q_2$  respectively, then  $\chi_1\chi_2$  is one modulo  $q_1q_2$ .
- If  $\chi$  is a character, then so is  $\bar{\chi}$ , and  $\chi\bar{\chi} = \bar{\chi}\chi = \chi_0$ .
- If  $\chi_1, \chi_2, \chi_3$  are characters modulo  $q$  and  $\chi_1\chi_2(x) = \chi_1\chi_3(x)$  for every  $x$ , then  $\chi_2 = \chi_3$ .

- If  $\chi_1$  and  $\chi_2$  are characters modulo  $q_1$  and  $q_2$  respectively, then  $\chi_1\chi_2$  is one modulo  $q_1q_2$ .
- If  $\chi$  is a character, then so is  $\bar{\chi}$ , and  $\chi\bar{\chi} = \bar{\chi}\chi = \chi_0$ .
- If  $\chi_1, \chi_2, \chi_3$  are characters modulo  $q$  and  $\chi_1\chi_2(x) = \chi_1\chi_3(x)$  for every  $x$ , then  $\chi_2 = \chi_3$ .
- Multiply by  $\bar{\chi}_1$ .



- Given  $x$  with  $(x, q) = 1$  and any character  $\chi_1$  modulo  $q$  we have

$$\sum_{\chi \pmod{q}} \chi(x) = \sum_{\chi \pmod{q}} \chi_1 \chi(x) = \chi_1(x) \sum_{\chi \pmod{q}} \chi(x).$$

- Given  $x$  with  $(x, q) = 1$  and any character  $\chi_1$  modulo  $q$  we have

$$\sum_{\chi \pmod{q}} \chi(x) = \sum_{\chi \pmod{q}} \chi_1 \chi(x) = \chi_1(x) \sum_{\chi \pmod{q}} \chi(x).$$

- Now we have the analogue of the previous lemma.

## Lemma 2

*If  $(x, q) = 1$  and there is a  $\chi_1$  such that  $\chi_1(x) \neq 1$ , then*

$$\sum_{\chi \pmod{q}} \chi(x) = 0.$$

*If there is no such  $\chi_1$ , then*

$$\sum_{\chi \pmod{q}} \chi(x) = h.$$

- Given  $x$  with  $(x, q) = 1$  and any character  $\chi_1$  modulo  $q$  we have

$$\sum_{\chi \pmod{q}} \chi(x) = \sum_{\chi \pmod{q}} \chi_1 \chi(x) = \chi_1(x) \sum_{\chi \pmod{q}} \chi(x).$$

- Now we have the analogue of the previous lemma.

## Lemma 2

*If  $(x, q) = 1$  and there is a  $\chi_1$  such that  $\chi_1(x) \neq 1$ , then*

$$\sum_{\chi \pmod{q}} \chi(x) = 0.$$

*If there is no such  $\chi_1$ , then*

$$\sum_{\chi \pmod{q}} \chi(x) = h.$$

- Can we always find such a  $\chi_1$  when  $x \not\equiv 1 \pmod{q}$ ?

- The answer is yes.

## Lemma 3

*Given  $x$  with  $(x, q) = 1$  and  $x \not\equiv 1 \pmod{q}$  there is a character  $\chi_1$  modulo  $q$  such that  $\chi_1(x) \neq 1$ .*

- The answer is yes.

## Lemma 3

*Given  $x$  with  $(x, q) = 1$  and  $x \not\equiv 1 \pmod{q}$  there is a character  $\chi_1$  modulo  $q$  such that  $\chi_1(x) \neq 1$ .*

- We give a quick and dirty proof. Since  $x \not\equiv 1 \pmod{q}$ , there is a prime power  $p^k$  such that  $p^k | q$  and  $p^k \nmid x - 1$ .

- The answer is yes.

### Lemma 3

*Given  $x$  with  $(x, q) = 1$  and  $x \not\equiv 1 \pmod{q}$  there is a character  $\chi_1$  modulo  $q$  such that  $\chi_1(x) \neq 1$ .*

- We give a quick and dirty proof. Since  $x \not\equiv 1 \pmod{q}$ , there is a prime power  $p^k$  such that  $p^k | q$  and  $p^k \nmid x - 1$ .
- If  $p$  is odd, or  $p = 2$  and  $k = 1$  or  $2$ , then we can choose a primitive root  $g$  modulo  $p^k$ . Then we define a character  $\chi_2(z; p^k)$  modulo  $p^k$  by taking

$$\chi_2(g^y; p^k) = e(y/\phi(p^k)).$$

- The answer is yes.

### Lemma 3

*Given  $x$  with  $(x, q) = 1$  and  $x \not\equiv 1 \pmod{q}$  there is a character  $\chi_1$  modulo  $q$  such that  $\chi_1(x) \neq 1$ .*

- We give a quick and dirty proof. Since  $x \not\equiv 1 \pmod{q}$ , there is a prime power  $p^k$  such that  $p^k | q$  and  $p^k \nmid x - 1$ .
- If  $p$  is odd, or  $p = 2$  and  $k = 1$  or  $2$ , then we can choose a primitive root  $g$  modulo  $p^k$ . Then we define a character  $\chi_2(z; p^k)$  modulo  $p^k$  by taking

$$\chi_2(g^y; p^k) = e(y/\phi(p^k)).$$

- Note that if  $g^y \not\equiv 1 \pmod{p^k}$ , then  $y \not\equiv 0 \pmod{\phi(p^k)}$ .

- The answer is yes.

### Lemma 3

*Given  $x$  with  $(x, q) = 1$  and  $x \not\equiv 1 \pmod{q}$  there is a character  $\chi_1$  modulo  $q$  such that  $\chi_1(x) \neq 1$ .*

- We give a quick and dirty proof. Since  $x \not\equiv 1 \pmod{q}$ , there is a prime power  $p^k$  such that  $p^k | q$  and  $p^k \nmid x - 1$ .
- If  $p$  is odd, or  $p = 2$  and  $k = 1$  or  $2$ , then we can choose a primitive root  $g$  modulo  $p^k$ . Then we define a character  $\chi_2(z; p^k)$  modulo  $p^k$  by taking

$$\chi_2(g^y; p^k) = e(y/\phi(p^k)).$$

- Note that if  $g^y \not\equiv 1 \pmod{p^k}$ , then  $y \not\equiv 0 \pmod{\phi(p^k)}$ .
- Now define

$$\chi_1(x) = \chi_2(x; p^k)\chi_0(x; qp^{-k})$$



- The answer is yes.

### Lemma 3

*Given  $x$  with  $(x, q) = 1$  and  $x \not\equiv 1 \pmod{q}$  there is a character  $\chi_1$  modulo  $q$  such that  $\chi_1(x) \neq 1$ .*

- We give a quick and dirty proof. Since  $x \not\equiv 1 \pmod{q}$ , there is a prime power  $p^k$  such that  $p^k | q$  and  $p^k \nmid x - 1$ .
- If  $p$  is odd, or  $p = 2$  and  $k = 1$  or  $2$ , then we can choose a primitive root  $g$  modulo  $p^k$ . Then we define a character  $\chi_2(z; p^k)$  modulo  $p^k$  by taking

$$\chi_2(g^y; p^k) = e(y/\phi(p^k)).$$

- Note that if  $g^y \not\equiv 1 \pmod{p^k}$ , then  $y \not\equiv 0 \pmod{\phi(p^k)}$ .
- Now define

$$\chi_1(x) = \chi_2(x; p^k)\chi_0(x; qp^{-k})$$

- That leaves the case when  $p = 2$  and  $k \geq 3$ , which is a little more complicated.

- Choose  $y, z$  so that

$$x \equiv (-1)^y 5^z \pmod{2^k}$$

- Choose  $y, z$  so that

$$x \equiv (-1)^y 5^z \pmod{2^k}$$

- Now we construct  $\chi_2$  as follows.

- Choose  $y, z$  so that

$$x \equiv (-1)^y 5^z \pmod{2^k}$$

- Now we construct  $\chi_2$  as follows.
- If  $y = 0$ , so that  $0 \leq z < 2^{k-2}$ , then take

$$\chi_2((-1)^u 5^v; 2^k) = e(v/2^{k-2}).$$

- Choose  $y, z$  so that

$$x \equiv (-1)^y 5^z \pmod{2^k}$$

- Now we construct  $\chi_2$  as follows.
- If  $y = 0$ , so that  $0 \leq z < 2^{k-2}$ , then take

$$\chi_2((-1)^u 5^v; 2^k) = e(v/2^{k-2}).$$

- If  $y = 1$ , then take

$$\chi_2((-1)^u 5^v; 2^k) = e(u/2).$$

- Choose  $y, z$  so that

$$x \equiv (-1)^y 5^z \pmod{2^k}$$

- Now we construct  $\chi_2$  as follows.
- If  $y = 0$ , so that  $0 \leq z < 2^{k-2}$ , then take

$$\chi_2((-1)^u 5^v; 2^k) = e(v/2^{k-2}).$$

- If  $y = 1$ , then take

$$\chi_2((-1)^u 5^v; 2^k) = e(u/2).$$

- Then proceed as before.

- We now can state the basic theorem for characters.

## Theorem 4

*There are  $\phi(q)$  characters modulo  $q$ ,*

$$\frac{1}{\phi(q)} \sum_{\chi \pmod{q}} \bar{\chi}(a)\chi(x) = \begin{cases} 1 & x \equiv a \pmod{q} \text{ \& } (a, q) = 1, \\ 0 & x \not\equiv a \pmod{q} \text{ or } (a, q) > 1. \end{cases}$$

*and*

$$\frac{1}{\phi(q)} \sum_{\chi \pmod{q}} \bar{\chi}_1(x)\chi_2(x) = \begin{cases} 1 & \chi_1 = \chi_2 \text{ and } (x, q) = 1, \\ 0 & \chi_1 \neq \chi_2. \end{cases}$$

- Consider the sum

$$\sum_{x \pmod{q}} \chi(x) \sum_{\chi \pmod{q}} \chi(x).$$



- Consider the sum

$$\sum_{x \pmod{q}} \sum_{\chi \pmod{q}} \chi(x).$$

- The sum over  $\chi$  contributes 0 if  $x \not\equiv 1 \pmod{q}$ ,  $h$  otherwise, so

$$= h.$$

- Interchanging the order gives

$$\begin{aligned} \sum_{\chi \pmod{q}} \sum_{x \pmod{q}} \chi(x) &= \sum_{x \pmod{q}} \chi_0(x) \\ &= \phi(q). \end{aligned}$$

- Given a character  $\chi$  modulo  $q$ , if there is a character  $\chi^*$  modulo  $r$ , with  $r|q$ , such that

$$\chi(x; q) = \chi^*(x; r)\chi_0(x; q),$$

then we say that  $\chi^*$  **induces**  $\chi$ .

- Given a character  $\chi$  modulo  $q$ , if there is a character  $\chi^*$  modulo  $r$ , with  $r|q$ , such that

$$\chi(x; q) = \chi^*(x; r)\chi_0(x; q),$$

then we say that  $\chi^*$  **induces**  $\chi$ .

- If there is no such character with  $r < q$ , then we say that  $\chi$  is **primitive**.

- Given a character  $\chi$  modulo  $q$ , if there is a character  $\chi^*$  modulo  $r$ , with  $r|q$ , such that

$$\chi(x; q) = \chi^*(x; r)\chi_0(x; q),$$

then we say that  $\chi^*$  **induces**  $\chi$ .

- If there is no such character with  $r < q$ , then we say that  $\chi$  is **primitive**.
- If  $\chi^*$  is primitive, then we call  $r$  the conductor of  $\chi$ .

- We now give two useful criteria for primitivity.

## Theorem 5

*Let  $\chi$  be a character modulo  $q$ . Then the following are equivalent:*

*(1)  $\chi$  is primitive.*

*(2) If  $d \mid q$  and  $d < q$  then there is a  $c$  such that  $c \equiv 1 \pmod{d}$ ,  $(c, q) = 1$ ,  $\chi(c) \neq 1$ .*

*(3) If  $d \mid q$  and  $d < q$ , then for every integer  $a$ ,*

$$\sum_{\substack{n=1 \\ n \equiv a \pmod{d}}}^q \chi(n) = 0.$$

- We now give two useful criteria for primitivity.

## Theorem 5

*Let  $\chi$  be a character modulo  $q$ . Then the following are equivalent:*

- (1)  $\chi$  is primitive.*
- (2) If  $d \mid q$  and  $d < q$  then there is a  $c$  such that  $c \equiv 1 \pmod{d}$ ,  $(c, q) = 1$ ,  $\chi(c) \neq 1$ .*
- (3) If  $d \mid q$  and  $d < q$ , then for every integer  $a$ ,*

$$\sum_{\substack{n=1 \\ n \equiv a \pmod{d}}}^q \chi(n) = 0.$$

- The proof is usually given in Math 568, and can be found in the files section.

- Given a character  $\chi$  modulo  $q$ , we define the Gauss sum  $\tau(\chi)$  of  $\chi$  to be

$$\tau(\chi) = \sum_{a=1}^q \chi(a)e(a/q).$$

- Given a character  $\chi$  modulo  $q$ , we define the Gauss sum  $\tau(\chi)$  of  $\chi$  to be

$$\tau(\chi) = \sum_{a=1}^q \chi(a)e(a/q).$$

- The Gauss sum is a special case of the more general sum

$$c_{\chi}(n) = \sum_{a=1}^q \chi(a)e(an/q).$$



- Given a character  $\chi$  modulo  $q$ , we define the Gauss sum  $\tau(\chi)$  of  $\chi$  to be

$$\tau(\chi) = \sum_{a=1}^q \chi(a)e(a/q).$$

- The Gauss sum is a special case of the more general sum

$$c_{\chi}(n) = \sum_{a=1}^q \chi(a)e(an/q).$$

- When  $\chi$  is the principal character, this is Ramanujan's sum

$$c_q(n) = \sum_{\substack{a=1 \\ (a,q)=1}}^q e(an/q),$$

- We now show that the sum  $c_\chi(n)$  is closely related to  $\tau(\chi)$ .

## Theorem 6

*Suppose that  $\chi$  is a character modulo  $q$ . If  $(n, q) = 1$  then*

$$\chi(n)\tau(\bar{\chi}) = \sum_{a=1}^q \bar{\chi}(a)e(an/q), \quad (1)$$

*and in particular*

$$\overline{\tau(\chi)} = \chi(-1)\tau(\bar{\chi}).$$

## Proof.

If  $(n, q) = 1$  then the map  $a \mapsto an$  permutes the residues modulo  $q$ , and hence

$$\chi(n)c_\chi(n) = \sum_{a=1}^q \chi(an)e(an/q) = \tau(\chi).$$

On replacing  $\chi$  by  $\bar{\chi}$ , this gives (6), and (7) follows by taking  $n = -1$ . □



- There is a multiplicative property of Gauss sums which is useful.

## Theorem 7

*Suppose that  $(q_1, q_2) = 1$ , that  $\chi_i$  is a character modulo  $q_i$  for  $i = 1, 2$ , and that  $\chi = \chi_1\chi_2$ . Then*

$$\tau(\chi) = \tau(\chi_1)\tau(\chi_2)\chi_1(q_2)\chi_2(q_1).$$

- There is a multiplicative property of Gauss sums which is useful.

## Theorem 7

*Suppose that  $(q_1, q_2) = 1$ , that  $\chi_i$  is a character modulo  $q_i$  for  $i = 1, 2$ , and that  $\chi = \chi_1\chi_2$ . Then*

$$\tau(\chi) = \tau(\chi_1)\tau(\chi_2)\chi_1(q_2)\chi_2(q_1).$$

- This is standard.

## Proof.

By the Chinese remainder theorem, each  $a \pmod{q_1q_2}$  can be written uniquely as  $a_1q_2 + a_2q_1$  with  $1 \leq a_i \leq q_i$ . Thus the general term in (3) is  $\chi_1(a_1q_2)\chi_2(a_2q_1)e(a_1/q_1) e(a_2/q_2)$ , so the result follows.  $\square$

- For primitive characters the hypothesis that  $(n, q) = 1$  in the first theorem can be removed.

## Theorem 8

*Suppose that  $\chi$  is a primitive character modulo  $q$ . Then*

$$\chi(n)\tau(\bar{\chi}) = \sum_{a=1}^q \bar{\chi}(a)e(an/q), \quad (2)$$

*holds for all  $n$ , and  $|\tau(\chi)| = \sqrt{q}$ .*

- We will make use of this when studying the large sieve.

## Proof.

It suffices to prove (2) when  $(n, q) > 1$ . Choose  $m$  and  $d$  so that  $(m, d) = 1$  and  $m/d = n/q$ . Then

$$\sum_{a=1}^q \chi(a) e(an/q) = \sum_{h=1}^d e(hm/d) \sum_{\substack{a=1 \\ a \equiv h \pmod{d}}}^q \chi(a).$$

Since  $d \mid q$  and  $d < q$ , the inner sum vanishes by Theorem 5. Thus (2) holds.  $\square$